Bellman-Ford Algorithm for Completion of Route Determination: An Experimental Study

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ABSTRACT

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Keywords:

Bellman-Ford Algorithm, Shortest route search, Traveling Salesman Problem, A review of Algorithms, Experimental study. In this study, a review of the existing Bellman-Ford Algorithm by conducting tests to see the accuracy of the route data or the shortest route. In this study, there are fifth locations that will be tested to see whether the route is really under the actual situation. The shortest path is part of the field of graph theory. If a graph has weight, then in the case of the shortest route, how can we do the minimization of the total weight of the route. This is what was done in this study to see how optimal the Bellman-Ford Algorithm is in handling the shortest route to be more accurate. The fifth Mall data is the most frequently visited by people in the city of Palembang. The five malls are Opi Mall, International Plaza, Palembang Indah Mall, Palembang Square, and Palembang Icon. The conclusion from the results of this study is that the Bellman-Ford Algorithm is more complicated to do in the search for calculations manually on the completion of the Traveling Salesman Problem (TSP). However, this algorithm is better in terms of finding optimal solutions and solving singe pair routes.

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1. INTRODUCTION

In the problems of daily life, finding the fastest route is still the priority of everyone, both going and traveling from one place to another. Nowadays, for every traveler, people always look for the fastest path to get to their destination immediately by utilizing digital maps. Digital maps utilize algorithm optimization so that the level of accuracy gets better. Some popular algorithms that can solve problems in finding the shortest route include: Dijkstra's Algorithm, Bellman-Ford Algorithm, Branch and Bound Algorithm, A-star Algorithm, Floyd-Warshall Algorithm, Ant Colony Algorithm, Vector Distance Algorithm, Ford-Fulkerson Algorithm [1]. The algorithm basically utilizes a graph model to facilitate the determination of nodes. The graph model is a complex and complicated mathematical model, but it can still be a solution that helps in various forms of a particular case. Graph models are beneficial in terms of modeling a problem [1] [2]. In the shortest route, the graph used a weighted graph from each side, is given a value or weighting [3]. In this study, the Bellman-Ford Algorithm is only used if there is a negative side [4]. So this study obtained two research objectives: to know the efficiency of the algorithm in the shortest route in terms of calculation and to test the algorithm in solving the shortest path problem.

In determining the shortest route from a source node, the Bellman-Ford Algorithm will calculate the shortest distance (from one particular vertex source) on a graph [1]. The purpose of one source is that this algorithm calculates all the shortest distances that start from one particular vertex point to all other vertices in a weighted and directed graph [5] [6]. Besides that, this algorithm uses d [u] as the upper limit with the distance d [u, v] from u to v. This algorithm makes improvements (updating) the distance at each source point to the point v in V until it reaches the Boolean TRUE theorem that the negative distance contains a non-negative

circle then the point can be reached from the source point [7]. In other conditions, it says Boolean FALSE [8] [9]. In the implementation, this algorithm also produces a path resulting from finding the shortest path [10]. Several applications of the Bellman-Ford Algorithm in a company location search application are to determine the company's destination [11], the user's starting point, the distance of the user's position to the position of the destination [12].

In this study, the focus of the use of the Bellman-Ford Algorithm is as part of the simulation application for locating the location of a mall in Palembang by looking at several points or the closest node to facilitate the identification process. Samples in the simulation use Mall data in Palembang, as many as 5 large locations in Palembang. The fifth Mall data is the most frequently visited by people in the city of Palembang. The five malls are OPI Mall, International Plaza, Palembang Indah Mall, Palembang Square, and Palembang Icon.

2. **RESEARCH METHOD**

The method used in the simulation process and application development uses a rational unified process (RUP). The iterative software development approach is repeated (iterative), focused on architecture (architecture-centric), more directed based on use cases (use case-driven) [13] [14]. While the research method used is the experimental method approach [15]. Research within the scope of the experiment is a laboratory, meaning that this research method is considered the most possible to conduct a hypothesis test between the causal relationship of fulfilling internal validity [16].

In this study, to conduct software analysis and design using the RUP development approach with the four phases owned as in Figure 1.



Fig 1. The lifecycle of the RUP Approach [13].

Inception a.

> This stage is more modeling the business processes needed (business modeling) and defining the need for a system to be made (requirements).

b. Elaboration

This stage is more focused on system architecture planning. This stage can also detect whether the desired system architecture can be made or not. Detecting the risks that might occur from the architecture created. This stage is more on system analysis and design and system implementation that focuses on the system prototype.

c. Construction

> This stage focuses on developing components and system features. This stage is more on the implementation and testing of systems that focus on implementing software in the program code. This stage produces a software product, a prerequisite for the Initial Operational Capability Milestone, or initial operational capability threshold.

d. Transition

> This stage is more on the deployment or installation of the system so that the user can understand it. This stage produces a software product, a prerequisite for the Initial Operational Capability Milestone, or an initial operational capability threshold. Activities at this stage include user training, system maintenance, and testing whether it meets user expectations.

Next, to conduct experiments with the RUP approach, the Bellman-Ford Algorithm is used to help the optimization process of determining the nearest path of the intended object [12]. The object of this research is the location of malls in Palembang. This algorithm must be following the procedure so that it can produce an optimal point, and no errors occur. The pseudocode of the algorithm is written as follows [17] [2]:

- 1. INITIALIZE-SINGLE-SOURCE (G, s)
- 2. for each vertex i = 1 to V(G) 1 do
- 3. for each edge (u, v) in E[G] do
- 4. RELAX (u, v, w)
- 5. For each edge (u, v) in E[G] do
- 6. if d[u] + w[u, v] < d[v] then

7. return FALSE

8. return TRUE.

This algorithm was chosen because it has the reliability in handling cases of negatively weighted graphs, can display alternative pathways, and the level of complexity of this algorithm is not too complicated [5] [17]. However, this algorithm also has a weakness, which is a graph that has a negative cycle [1][2], so for graphs that have negative cycles, it cannot calculate the shortest path. Below is the flow of the Bellman-Ford algorithm procedure as in algorithm 1 [1].

Data: G, s

Result: shortest paths between s and any other vertex of G; if there exists a negative cycle, it outputs its existence

 $\begin{array}{ll} 1. \mbox{ for } u \in V \mbox{ do} \\ 2. & l(u): = +\infty; \\ 3. \mbox{ l(s): = 0;} \\ 4. \mbox{ for } I = 1 \mbox{ to } |V| - 1 \mbox{ do} \\ 5. & \mbox{ foreach } edge \ \{u,v\} + w(\{u,v\}); \\ 6. & l(u): = \min((l(u), l(v) + w(\{u,v\}); \\ 7. \mbox{ foreach } edge \ \{u,v\} \in E \mbox{ do} \\ 8. & \mbox{ if } l(u) > l(v) + w(\{u,v\}) \mbox{ then} \\ 9. & \mbox{ Return ("Negative cycle");} \end{array}$

Script 1. Bellman-Ford Shortest Path Algorithm [1]

2.1. Design

The design process carried out in this simulation involves the admin and user in every interaction. Users can use it by accessing using a mobile web device or browser, as in Figure 2.



Fig 2. User interaction with the system

Next, to see the activities between the admin and the user, an activity diagram is designed. Each sequence of activities described is a defined business process system, sequence, or grouping of views from the interface. Each activity is considered to have a display interface design, test design. Each activity is considered to require a test that needs to be defined test cases and menu designs displayed on the software. The following activities occur in the application shown in Figure 3.





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2.1. Data Sample

This study used 5 large locations Mall data in Palembang (Table 1) as samples. The Mall data is the most frequently visited by people in the city of Palembang. The five malls are OPI Mall, International Plaza, Palembang Indah Mall, Palembang Square, and Palembang Icon. The starting point (node) in this simulation is Bina Darma University (UBD), which is located at Ahmad Yani Street No.3 Silaberanti Village, Seberang Ulu 1, Palembang.

Table 1. Data Sample starting point from UBD				
No.	Destination Mall	Mileage (Meters)		
1	International Plaza	1500		
2	Palembang Indah Mall	1610		
3	Palembang Icon	2340		
4	Palembang Square	2660		
5	OPI Mall	2920		

3. RESULTS AND DISCUSSION

In this discussion, researchers conducted a test using 1 Mall data from 5 existing data. The reason is because of the limitations of the pages that are allowed, so if of the fifth data discussed, there will certainly not be enough pages. However, for testing the fifth route of the mall data, we will present it in Table 4. The data to be tested is the route from UBD to OPI Mall with 2 alternative routes, where the starting point is symbolized by the letter S, and the other points are symbolized by numbers 2-13. The following are two alternative path test data, the calculation process of the Bellman-Ford Algorithm, and the interactions of each visit shown in Table 2 and Table 3 and Figure 4.

 Table 2. Alternative Pathways 1 and Alternative 2 testing data

Starting Point	Alternative Pathway 1	Alternative Pathway 2
UBD	Flyover Jakabaring	Flyover Jakabaring
	Lrt Kejaksaan Station	Lrt Kejaksaan Station
	Pangeran Ratu Street	Pangeran Ratu Street
	Pasar Induk market	LRT Jakababaring
		Station
	Yusnawati tailor	Warung Nasi Chintya
	Depot Kusen, Pipa Raya	DJKA Palembang
	Street	
	Alfamart Opi Raya	Opi Mall
	Perumnas GGS	
	Opi Mall	



Fig 4. Stage of grading each node



Based on the results, the route from point 1 to point 13, there are three iterations, which are: 1. First iteration

Iteration in Table 3, number 1, is iterations from point 1 to point 2, 3, and 4. In the first iteration, point 2 is filled with the distance from point 1 to point 2; point 3 is filled with distance from point 2 to point 3; point 4 is filled with the distance from point 3 to point 4. The minimum distance obtained from point 1 to point 2 is 0 + 200 = 200 m, point 2 to point 3 is 200 + 480 = 680 m and point 3 to point 4 is 680 + 830 = 1510 m.

2. Second iteration

The following are iterations in Table 3 number 2. Point 5 is filled with the distance from point 4 to point 5. Point 6 is filled with the distance from point 5 to point 6. Point 7 is filled with the distance from point 6 to point 7. Point 8 is filled with distance from point 7 to point 8. Point 9 is filled with the distance from point 8 to point 9. Point 13 is filled with the distance from point 9 to point 13.

Meanwhile, the minimum distance obtained from point 4 to point 5 is 1510 + 2110 = 3620 m; point 5 to point 6 is 3620 + 690 = 4310 m; point 6 to point 7 is 4310 + 880 = 5190 m; point 7 to point 8 is 5190 + 470 = 5660 m; point 8 to point 9 is 5660 + 530 = 6190 m; and point 9 to point 13 is 6190 + 300 = 6490 m.

3. Third Iteration

The iteration in Table 3, number 2, is the first iteration of point 10 filled with the distance from point 4 to point 10, point 11 is filled with the distance from point 10 to point 11, point 12 is filled with distance from point 11 to point 12, and point 13 is filled with the distance from point 12 to point 13. The minimum distance obtained from point 4 to point 10 is 1510 + 1880 = 3390 m, point 10 to point 11 is 3390 + 370 = 3760 m, point 11 to point 12 is 3760 + 910 = 4670 m, point 12 to point 13 is 4670 + 390 = 5060 m.

The following is the recapitulation results from the above calculations:

1. The first route is 1-2-3-4-5-6-7-8-9-13, with a distance of 6490 m / (6.49 km).

2. The second route is 1-2-3-4-10-11-12-13, with a distance of 5060 m / (5.06 km).

Table 4. The results of testing the five mall routes						
No	Starting Point	Destination	Mileage	Pathway		
			(Meters)			
1	UBD	International	3390	UBD→Flyover Jakabaring→		
		Plaza Mall		Bundunds of Air Mancur \rightarrow Masjid		
		(IP)		Agung \rightarrow Puspita Photo Printing \rightarrow		
				Agung Shop Sport \rightarrow IP.		
2	UBD	OPI Mall	5060	$UBD \rightarrow Flyover Jakabaring \rightarrow LRT$		
				Kejaksaan Station \rightarrow Pangeran Ratu		
				Street \rightarrow LRT Jakabaring		
				Station→Warung Nasi		
				Chintya \rightarrow DJKA Palembang \rightarrow OPI		
				Mall.		
3	UBD	Palembang	3420	UBD→Flyover Jakabaring→		
		Indah Mall		Bundunds of Air Mancur \rightarrow Mayor's		
		(PIM)		office \rightarrow Hospital for Lungs \rightarrow South		
		D 1 1	11.00	Station \rightarrow Red Light PIM \rightarrow PIM.		
4	UBD	Palembang Icon Mall (PI)	4130	UBD→Flyover Jakabaring→		
				Bundunds of Air Mancur → Puspita		
				Photo Printing→Agung Shop Sport→		
				Agate Sales Center Cinde Market		
				(cinde) → Intersection 4 Charitas		
				Hospital \rightarrow Dinas Pernubungan Office		
				→ Palembang emas Darusalam		
5		Dolombong	4770	Park 7P1		
5	UBD	Squara Mall	4770	Bundunds of Air Mongur Duspite		
		(DS)		Duridunds of All Mancul γ fuspita Dhoto Drinting $\rightarrow \Lambda$ gung Shop Sport \rightarrow		
		(13)		A gata Salas Contar Cinda Markat		
				$(\text{cinde}) \rightarrow \text{Intersection } A \text{ Charitas}$		
				Hospital \rightarrow Dinas Parbubungan		
				Office -> Palambang amas Darusalam		
				$Park \rightarrow PS$		
				raik /r.s		

4. CONCLUSION

The results of this study, the Bellman-Ford Algorithm is more complicated to do in the search for calculations manually on the completion of the Traveling Salesman Problem (TSP). However, this algorithm is better in terms of finding optimal solutions and solving singe pair routes.

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